# On the Gaussian surface area of spectrahedra 

Srinivasan Arunachalam* Oded Regev ${ }^{\dagger}$ Penghui $\mathrm{Yao}^{\ddagger}$

December 2, 2021


#### Abstract

We show that for sufficiently large $n \geq 1$ and $d=C n^{3 / 4}$ for some universal constant $C>0$, a random spectrahedron with matrices drawn from Gaussian orthogonal ensemble has Gaussian surface area $\Theta\left(n^{1 / 8}\right)$ with high probability.


## 1 Introduction

A spectrahedron $S \subseteq \mathbb{R}^{n}$ is a set of the form

$$
S=\left\{x \in \mathbb{R}^{n}: \sum_{i} x_{i} A^{(i)} \preceq B\right\},
$$

for some $d \times d$ symmetric matrices $A^{(1)}, \ldots, A^{(n)}, B \in \operatorname{Sym}_{d}$. Here we will be concerned with the Gaussian surface area of $S$, defined as

$$
\begin{equation*}
\operatorname{GSA}(S)=\liminf _{\delta \rightarrow 0} \frac{\mathcal{G}^{n}\left(S_{\delta}^{\text {out }}\right)}{\delta}, \tag{1}
\end{equation*}
$$

where $S_{\delta}^{\text {out }}=\{x \notin S: \operatorname{dist}(x, S) \leq \delta\}$ denotes the outer $\delta$-neighborhood of $S$ under Euclidean distance and $\mathcal{G}^{n}(\cdot)$ denotes the standard Gaussian measure on $\mathbb{R}^{n}$. Ball showed that the GSA of any convex body in $\mathbb{R}^{n}$ is $O\left(n^{1 / 4}\right)$ [Bal93], which was later shown to be tight by Nazarov [Naz03]. Moreover, Nazarov [KOS08] showed that the GSA of a $d$-facet polytope ${ }^{1}$ in $\mathbb{R}^{n}$ is $O(\sqrt{\log d})$ and this fact has found application in constructing pseudorandom generators for polytopes [HKM13, ST17, CDS19]. Motivated by recent work [AY21], this raises the question of whether the GSA of spectrahedra is also small. In this note we answer this question in the negative.

Theorem 1. For a universal constant $C>0$ and any integers $n, d \geq 1$ satisfying $d \leq n / C$ the following hold. If $\boldsymbol{A}^{(1)}, \ldots, \boldsymbol{A}^{(n)}$ are i.i.d. drawn from the $d \times d$ Gaussian orthogonal ensemble (see below for the definition), then the spectrahedron

$$
\begin{equation*}
\mathcal{T}=\left\{x \in \mathbb{R}^{n}: \sum_{i} x_{i} \boldsymbol{A}^{(i)} \preceq 2 \sqrt{n d} \cdot \mathbb{I}\right\} \tag{2}
\end{equation*}
$$

satisfies $\operatorname{GSA}(\mathcal{T})=\Omega(\sqrt{n / d})$ with probability at least $1-C \exp \left(-d n^{-3 / 4} / C\right)$. Moreover, for any integer d satisfying $d \leq n / C, \operatorname{GSA}(\mathcal{T})=O(\sqrt{n / d})$ holds with probability at least $1-\exp (-n / 50)$.

[^0]The theorem shows the existence of spectrahedra with GSA of $\Omega\left(n^{1 / 8}\right)$. (In fact, a random spectrahedron as above satisfies this with constant probability). This lower bound is close to the GSA upper bound of Ball [Bal93] of $O\left(n^{1 / 4}\right)$ for arbitrary convex bodies. Moreover, the lower bound shows that in contrast to the case of polytopes, the GSA of spectrahedra can depend polynomially on d. A natural open question is how large the GSA of arbitrary spectrahedra can be; can spectrahedra achieve a GSA of $\Theta\left(n^{1 / 4}\right)$ ?

## 2 Preliminaries

We use $\boldsymbol{g}, \boldsymbol{x}, \boldsymbol{A}$ to denote random variables. We let $\mathcal{G}\left(0, \sigma^{2}\right)$ be the normal distribution with mean 0 and variance $\sigma^{2}$. We denote by $\mathcal{H}_{d}$ the $d \times d$ Gaussian orthogonal ensemble (GOE). Namely, $\boldsymbol{A} \sim \mathcal{H}_{d}$ if it is a symmetric matrix with $\left\{\boldsymbol{A}_{i, j}\right\}_{i \leq j}$ independently distributed satisfying $\boldsymbol{A}_{i, j} \sim \mathcal{G}(0,1)$ for $i<j$ and $\boldsymbol{A}_{i, i} \sim \mathcal{G}(0,2)$. To keep notations short, for $b \geq 0$ we use $[a \pm b]$ to represent the interval $[a-b, a+b]$. For every $c \geq 0$, we use $c \cdot[a \pm b]$ to represent the interval $[a c \pm b c]$. We denote the set of $n$-dimensional unit vectors by $S^{n-1}$. Finally, we let $\chi_{n}$ be the $\chi$ distribution with $n$ degrees of freedom, which is the square root of the sum of the squares of $n$ independent standard normal variables. The following are some simple facts about the $\chi$ distribution.

Fact 2. Let $n \in \mathbb{Z}_{>0}$ and $h(\cdot)$ be the pdf of $\chi_{n}$. Then $h(x) \geq c$ for $x \in[\sqrt{n} \pm c]$, where $c>0$ is an absolute constant.

Fact 3. Let $n \in \mathbb{Z}_{>0}$ and $h(\cdot)$ be the pdf of $\chi_{n}$. Then $h(x) \leq O(\sqrt{n} /|x|)$ for $x \in \mathbb{R}$.
Proof. Recall that by definition

$$
h(x)=\frac{1}{2^{\frac{n}{2}-1} \Gamma\left(\frac{n}{2}\right)} x^{n-1} e^{-x^{2} / 2}
$$

for $x \geq 0$, and $h(x)=0$ otherwise. Hence the fact is trivial for $x \leq 0$. For $x>0$, the fact follows from the inequalities $x^{n} e^{-x^{2} / 2} \leq n^{n / 2} e^{-n / 2}$ and $\Gamma(z) \geq \sqrt{2 \pi} z^{z-1 / 2} e^{-z}$ for all $z>0$ [AAR99, Jam15].

Lemma 4 ([LM00, comment below Lemma 1]). For $n \geq 1$, let $\boldsymbol{r}$ be a random variable distributed according to $\chi_{n}$. Then for every $x>0$, we have

$$
\operatorname{Pr}\left[n-2 \sqrt{n x} \leq \boldsymbol{r}^{2} \leq n+2 \sqrt{n x}+2 x\right] \geq 1-2 e^{-x}
$$

For our purposes, it will be convenient to use an alternative definition of Gaussian surface area in terms of the inner surface area. Namely, for $S_{\delta}^{\text {in }}=\left\{x \in S: \operatorname{dist}\left(x, S^{c}\right\} \leq \delta\right)$ where $S^{c}$ is the complement of the body $S$, we define,

$$
\begin{equation*}
\operatorname{GSA}(S)=\lim _{\delta \rightarrow 0} \frac{\mathcal{G}^{n}\left(S_{\delta}^{\mathrm{in}}\right)}{\delta} \tag{3}
\end{equation*}
$$

It follows from Huang et al. [HXZ21, Theorem 3.3] that this definition is equal to the one in Eq. (1) when $S$ is a convex body that contains the origin, which is sufficient for our purposes.

To prove our main theorem, we use the following facts, starting with a well known bound on the size of an $\varepsilon$-net of the $n$-dimensional sphere.

Fact 5 ([Tao12, Lemma 2.3.4]). For every $d \geq 1$ and any $0<\varepsilon<1$ there exists an $\varepsilon$-net of the sphere $S^{d-1}$ of cardinality at most $(C / \varepsilon)^{d}$ for some universal constant $C>0$.

Claim 6 ([RS15, Page 134, Theorem 3]). Let $\boldsymbol{x}, \boldsymbol{y}$ be two real-valued random variables and $f$ be the pdf of $(\boldsymbol{x}, \boldsymbol{y})$. Then the pdf of $\boldsymbol{z}=\boldsymbol{x} \cdot \boldsymbol{y}$ is given by

$$
g(z)=\int_{-\infty}^{\infty} f\left(x, \frac{z}{x}\right) \cdot \frac{1}{|x|} d x .
$$

Theorem 7 ([LR10, Theorem 1]). Let $\boldsymbol{A} \sim \mathcal{H}_{d}$. For every $0<\eta<1$, it holds that

$$
\operatorname{Pr}\left[\lambda_{\max }(\boldsymbol{A}) \in 2 \sqrt{d}[1 \pm \eta]\right] \geq 1-C \cdot e^{-d \eta^{3 / 2} / C}
$$

for some absolute constant $C>0$.

## 3 Proof of main theorem

The core of the argument is in the following lemma, bounding $q(2 \sqrt{n d})$ where $q$ is the pdf of the largest eigenvalue of the matrix showing up in Eq. (2). We will later show that this value is essentially the same as $\operatorname{GSA}(\mathcal{T})$, where $\mathcal{T}$ is the spectrahedron in the statement of the theorem.

Lemma 8. For $n, d \geq 1$ and $A^{(1)}, \ldots, A^{(n)} \in \operatorname{Sym}_{d}$, let $q(\cdot)$ be the probability density function of

$$
\lambda_{\max }\left(\sum_{i} \boldsymbol{x}_{i} A^{(i)}\right)
$$

where $\boldsymbol{x}=\left(\boldsymbol{x}_{1}, \ldots, \boldsymbol{x}_{n}\right)$ is a random vector and each entry is i.i.d. drawn from $\mathcal{G}(0,1)$. If $\boldsymbol{A}^{(1)}, \ldots, \boldsymbol{A}^{(n)}$ are i.i.d. drawn from the $d \times d$ Gaussian orthogonal ensemble, then $q(2 \sqrt{n d})=\Omega(\sqrt{1 / d})$ with probability at least $1-C \exp \left(-d n^{-3 / 4} / C\right)$ (over the choice of $\boldsymbol{A}^{(1)}, \ldots, \boldsymbol{A}^{(n)}$ ) where $C>0$ is a universal constant. Moreover, for any integer $d$ and any $d \times d$ matrices $A^{(1)}, \ldots, A^{(n)}, q(2 \sqrt{n d})=O(\sqrt{1 / d})$.

Proof. Let $\boldsymbol{y} \sim S^{n-1}$ be chosen uniformly from the unit sphere and for matrices $A^{(1)}, \ldots, A^{(n)}$, denote by $p$ the pdf of $\lambda_{\max }\left(\sum_{i} \boldsymbol{y}_{i} A^{(i)}\right)$. Let $\boldsymbol{r} \sim \chi_{n}$ and notice that $\boldsymbol{r} \boldsymbol{y}$ is distributed like $\boldsymbol{x}$. Denote by $h$ the pdf of $\boldsymbol{r}$. By Claim 6, we have

$$
\begin{equation*}
q(2 \sqrt{n d})=\int_{-\infty}^{\infty} h(2 \sqrt{n d} / z) p(z) \frac{1}{|z|} d z \tag{4}
\end{equation*}
$$

Using Fact $3, h(2 \sqrt{n d} / z) /|z|=O(1 / \sqrt{d})$ for all $z$. Hence Eq. (4) can be bounded as $O(1 / \sqrt{d}) \int_{-\infty}^{\infty} p(z) d z=$ $O(1 / \sqrt{d})$, establishing the claimed upper bound on $q$.

To prove the lower bound on $q$, let $\boldsymbol{A}^{(1)}, \ldots, \boldsymbol{A}^{(n)} \sim \mathcal{H}_{d}$ be $n$ matrices chosen i.i.d. from the GOE. Observe that by Theorem 7, we have

$$
\begin{equation*}
\operatorname{Pr}\left[\lambda_{\max }\left(\sum_{i=1}^{n} \boldsymbol{y}_{i} \boldsymbol{A}^{(i)}\right) \in I\right] \geq 1-C \exp \left(-d n^{-3 / 4} / C\right) \tag{5}
\end{equation*}
$$

where

$$
I=2 \sqrt{d} \cdot[1 \pm c / \sqrt{n}]
$$

for some universal constants $C, c>0$. Define the set of matrices

$$
G=\left\{\left(A^{(1)}, \ldots, A^{(n)}\right): \operatorname{Pr}\left[\lambda_{\max }\left(\sum_{i=1}^{n} \boldsymbol{y}_{i} A^{(i)}\right) \in I\right] \geq \frac{1}{2}\right\} .
$$

Then, using the definition of $G$ and Eq. (5), we have

$$
\operatorname{Pr}\left[\left(\boldsymbol{A}^{(1)}, \ldots, \boldsymbol{A}^{(n)}\right) \in G\right] \geq 1-2 C \exp \left(-d n^{-3 / 4} / C\right)
$$

Now fix any $\left(A^{(1)}, \ldots, A^{(n)}\right) \in G$. By definition of $G, \int_{I} p(z) d z \geq 1 / 2$, and therefore the right-hand side of Eq. (4) is at least

$$
\begin{equation*}
\int_{I} h(2 \sqrt{n d} / z) p(z) \frac{1}{z} d z \geq \Omega(1) \cdot \int_{I} p(z) \frac{1}{z} d z \geq \Omega(1 / \sqrt{d}) \tag{6}
\end{equation*}
$$

where we used Fact 2 to conclude that $h(2 \sqrt{n d} / z) \geq \Omega(1)$ for all $z \in I$.
We next relate $q(2 \sqrt{n d})$ to $\operatorname{GSA}(\mathcal{T})$. For a vector $v \in S^{d-1}$, and $d \times d$ symmetric matrices $A^{(1)}, \ldots, A^{(n)}$, define the vector

$$
\begin{equation*}
W_{v}=\left(v^{T} A^{(1)} v, v^{T} A^{(2)} v, \ldots, v^{T} A^{(n)} v\right) \in \mathbb{R}^{n} . \tag{7}
\end{equation*}
$$

Notice that $\mathcal{T}$ can be written as

$$
\mathcal{T}=\left\{x \in \mathbb{R}^{n}: \sum_{i} x_{i} A^{(i)} \preceq 2 \sqrt{n d} \cdot \mathbb{I}\right\}=\left\{x \in \mathbb{R}^{n}: \forall v \in S^{d-1},\left\langle x, W_{v}\right\rangle \leq 2 \sqrt{n d}\right\} .
$$

We say that $A^{(1)}, \ldots, A^{(n)}$ are good if

$$
\forall v \in S^{d-1}, \frac{1}{2} \sqrt{n} \leq\left\|W_{v}\right\| \leq 2 \sqrt{n}
$$

Lemma 9. There exists a constant $C \geq 1$ such that for all integers $n$ and $d \leq n / C$, random matrices $\boldsymbol{A}^{(1)}, \ldots, \boldsymbol{A}^{(n)}$ drawn i.i.d. from $\mathcal{H}_{d}$ are good with probability at least $1-\exp (-n / 50)$.

Proof. For a fixed $v \in S^{d-1}$, we claim that

$$
\begin{equation*}
\operatorname{Pr}\left[n \leq\left\|\boldsymbol{W}_{v}\right\|^{2} \leq 3 n\right] \geq 1-2 \exp (-n / 40) . \tag{8}
\end{equation*}
$$

To see this, observe that for $\boldsymbol{A} \sim \mathcal{H}_{d}$ and unit vector $v \in \mathbb{R}^{d}, v^{T} \boldsymbol{A} v$ is distributed according to

$$
\left(4 \sum_{i<j} v_{i}^{2} v_{j}^{2}+2 \sum_{i} v_{i}^{4}\right)^{1 / 2} \cdot \mathcal{G}(0,1)=\sqrt{2} \cdot \mathcal{G}(0,1)
$$

Therefore, each entry in $\boldsymbol{W}_{v}$ is distributed according to $\mathcal{G}(0,2)$, and Lemma 4 implies Eq. (8). We next prove that with high probability (over the $\boldsymbol{A}^{(i)} \mathrm{s}$ ), for every unit vector $z,\left\|\boldsymbol{W}_{z}\right\|$ is large. First, by Fact 5 , there exists a set $\mathcal{V}=\left\{v_{1}, \ldots, v_{\left(10^{4} M\right)}\right\} \subseteq \mathbb{R}^{d}$ of unit vectors that form a $10^{-4}$-net of the unit Euclidean sphere where $M$ is a constant. Applying a union bound on $\mathcal{V}$, we have

$$
\begin{equation*}
\operatorname{Pr}\left[\forall v \in \mathcal{V}: n \leq\left\|\boldsymbol{W}_{v}\right\|^{2} \leq 3 n\right] \geq 1-2 \exp (-n / 40) \cdot\left(10^{4} M\right)^{d} \geq 1-\exp (-n / 50), \tag{9}
\end{equation*}
$$

here we used that $M$ is a constant and $d \leq n / C$ for a sufficiently large $C$.
To conclude the proof, it suffices to show that if $\boldsymbol{A}^{(1)}, \ldots, \boldsymbol{A}^{(n)}$ are such that

$$
\forall v \in \mathcal{V}, n \leq\left\|\boldsymbol{W}_{v}\right\|^{2} \leq 3 n
$$

then also

$$
\forall z \in S^{d-1},\left\|\boldsymbol{W}_{z}\right\| \geq \frac{1}{2} \sqrt{n}
$$

Let $\boldsymbol{b}_{\max }=\max _{z \in S^{d-1}}\left\|\boldsymbol{W}_{z}\right\|$ and $\boldsymbol{b}_{\min }=\min _{z \in S^{d-1}}\left\|\boldsymbol{W}_{z}\right\|$. Let $\boldsymbol{z}_{\max }$ and $\boldsymbol{z}_{\min }$ be the vectors achieving the maximum and the minimum respectively. Let $\boldsymbol{v}_{\max }$ and $\boldsymbol{v}_{\min }$ be the vectors in $\mathcal{V}$ that are closest to $\boldsymbol{z}_{\max }$ and $\boldsymbol{z}_{\min }$, respectively. For any vectors $z, v \in S^{d-1}$ with $\|z-v\| \leq 10^{-4}$, applying the spectral decomposition of $z z^{T}-v v^{T}$, there exist unit vectors $u_{1}, u_{2}$ and $0 \leq \lambda \leq \frac{1}{100}$ such that

$$
\begin{equation*}
z z^{T}-v v^{T}=\lambda \cdot\left(u_{1} u_{1}^{T}-u_{2} u_{2}^{T}\right) \tag{10}
\end{equation*}
$$

Hence

$$
\begin{aligned}
\left\|\boldsymbol{W}_{z}-\boldsymbol{W}_{v}\right\|^{2}=\sum_{i=1}^{n}\left(z^{T} \boldsymbol{A}^{(i)} z-v^{T} \boldsymbol{A}^{(i)} v\right)^{2} & =\sum_{i=1}^{n}\left(\operatorname{Tr}\left(\boldsymbol{A}^{(i)}\left(z z^{T}-v v^{T}\right)\right)\right)^{2} \\
& \leq \frac{1}{10^{4}} \sum_{i=1}^{n}\left(u_{1}^{T} \boldsymbol{A}^{(i)} u_{1}-u_{2}^{T} \boldsymbol{A}^{(i)} u_{2}\right)^{2} \\
& \leq \frac{1}{5000} \sum_{i=1}^{n}\left(\left(u_{1}^{T} \boldsymbol{A}^{(i)} u_{1}\right)^{2}+\left(u_{2}^{T} \boldsymbol{A}^{(i)} u_{2}\right)^{2}\right) \\
& \leq \frac{\boldsymbol{b}_{\max }^{2}}{2500}
\end{aligned}
$$

Choosing $z=\boldsymbol{z}_{\max }$ and $v=\boldsymbol{v}_{\max }$, we have

$$
\left\|\boldsymbol{W}_{\boldsymbol{z}_{\max }}\right\| \leq\left\|\boldsymbol{W}_{\boldsymbol{v}_{\max }}\right\|+\frac{\boldsymbol{b}_{\max }}{50}
$$

Now, since $\left\|\boldsymbol{W}_{\boldsymbol{z}_{\max }}\right\|=\boldsymbol{b}_{\max }$, we have

$$
\boldsymbol{b}_{\max } \leq \frac{50}{49}\left\|\boldsymbol{W}_{\boldsymbol{v}_{\max }}\right\| \leq \frac{50}{49} \sqrt{3 n} \leq 2 \sqrt{n}
$$

Similarly, we set $z=\boldsymbol{z}_{\text {min }}$ and $v=\boldsymbol{v}_{\text {min }}$ and obtain

$$
\boldsymbol{b}_{\min } \geq\left\|\boldsymbol{W}_{\boldsymbol{v}_{\min }}\right\|-\frac{\boldsymbol{b}_{\max }}{50} \geq \sqrt{n}-\frac{1}{25} \sqrt{n}>\frac{1}{2} \sqrt{n}
$$

This concludes the result.
For the following claim, we define the inner and outer shells of $\mathcal{T}$ as

$$
\begin{aligned}
\mathcal{D}_{\delta}^{\text {in }} & =\left\{x: \lambda_{\max }\left(\sum_{i} x_{i} A^{(i)}\right) \in \sqrt{n} \cdot[2 \sqrt{d}-\delta, 2 \sqrt{d}]\right\} \\
\mathcal{D}_{\delta}^{\text {out }} & =\left\{x: \lambda_{\max }\left(\sum_{i} x_{i} A^{(i)}\right) \in \sqrt{n} \cdot[2 \sqrt{d}, 2 \sqrt{d}+\delta]\right\}
\end{aligned}
$$

Also recall the inner and outer neighborhoods of $\mathcal{T}$, defined as

$$
\begin{aligned}
\mathcal{T}_{\delta}^{\text {in }} & =\{x \in \mathcal{T}: \exists y \notin \mathcal{T}:\|x-y\| \leq \delta\} \\
\mathcal{T}_{\delta}^{\text {out }} & =\{x \notin \mathcal{T}: \exists y \in \mathcal{T}:\|x-y\| \leq \delta\}
\end{aligned}
$$

Claim 10. For sufficiently small $\delta>0$ and any good $A^{(1)}, \ldots, A^{(n)}$, we have $\mathcal{D}_{\delta}^{\mathrm{in}} \subseteq \mathcal{T}_{4 \delta}^{\mathrm{in}}$ and $\mathcal{T}_{\delta}^{\text {out }} \subseteq \mathcal{D}_{2 \delta}^{\text {out }}$.

Proof. For every $x \in \mathcal{D}_{\delta}^{\text {in }}$, let $v$ be a unit eigenvector of $\sum_{i} x_{i} A^{(i)}$ with the eigenvalue $\lambda_{\max }\left(\sum_{i} x_{i} A^{(i)}\right)$. Therefore,

$$
\left\langle x, W_{v}\right\rangle=v^{T}\left(\sum x_{i} A^{(i)}\right) v \geq(2 \sqrt{d}-\delta) \sqrt{n} .
$$

Setting $y=2 \delta \sqrt{n} W_{v} /\left\|W_{v}\right\|^{2}$, we have

$$
\left\langle x+y, W_{v}\right\rangle=\left\langle x, W_{v}\right\rangle+2 \delta \sqrt{n} \geq(2 \sqrt{d}-\delta) \sqrt{n}+2 \delta \sqrt{n}=(2 \sqrt{d}+\delta) \sqrt{n}
$$

and so $x+y \notin \mathcal{T}$. Moreover, since $A^{(1)}, \ldots, A^{(n)}$ are good, $\|y\|=2 \delta \sqrt{n} /\left\|W_{v}\right\| \leq 4 \delta$ and therefore $x \in \mathcal{T}_{4 \delta}^{\text {in }}$, as desired. For the other containment, let $x \in \mathcal{T}_{\delta}^{\text {out }}$. Then for any unit vector $v$, by Cauchy-Schwarz and using $\left\|W_{v}\right\| \leq 2 \sqrt{n}$,

$$
\left\langle x, W_{v}\right\rangle \leq 2 \sqrt{n d}+2 \delta \sqrt{n}
$$

implying that $x \in \mathcal{D}_{2 \delta}^{\text {out }}$, as desired.
We now prove our main theorem.
Proof of Theorem 1. First observe that since $q(\cdot)$ is continuous, the lower bound on the pdf in Lemma 8 implies that $\mathcal{G}^{n}\left(\mathcal{D}_{\delta}^{\text {in }}\right) \geq \Omega(\delta \sqrt{n / d})$ for sufficiently small $\delta>0$. Thus, $\mathcal{G}^{n}\left(\mathcal{T}_{4 \delta}\right)=\Omega(\delta \sqrt{n / d})$ by Claim 10. By definition of $\operatorname{GSA}(S)=\lim _{\delta \rightarrow 0} \mathcal{G}^{n}\left(S_{\delta}^{\text {in }}\right) / \delta$, we obtain the desired lower bound on GSA. Similarly, using the upper bound on the pdf in Lemma $8, \mathcal{G}^{n}\left(\mathcal{D}_{\delta}^{\text {out }}\right)=O(\delta \sqrt{n / d})$ for sufficiently small $\delta>0$. Thus, $\mathcal{G}^{n}\left(\mathcal{T}_{\delta / 2}^{\text {out }}\right)=O(\delta \sqrt{n / d})$ by Claim 10 . We complete the proof using $\operatorname{GSA}(S)=\lim _{\delta \rightarrow 0} \mathcal{G}^{n}\left(S_{\delta}^{\text {out }}\right) / \delta$.

Acknowledgements. We thank Daniel Kane, Assaf Naor, Fedor Nazarov, and Yiming Zhao for useful correspondence. O.R. is supported by the Simons Collaboration on Algorithms and Geometry, a Simons Investigator Award, and by the National Science Foundation (NSF) under Grant No. CCF-1814524. P.Y. is supported by the National Key R\&D Program of China 2018YFB1003202, National Natural Science Foundation of China (Grant No. 61972191), the Program for Innovative Talents and Entrepreneur in Jiangsu and Anhui Initiative in Quantum Information Technologies Grant No. AHY150100.

## References

[AAR99] George E. Andrews, Richard Askey, and Ranjan Roy. The Gamma and Beta Functions, page 1-60. Encyclopedia of Mathematics and its Applications. Cambridge University Press, 1999.
[AY21] Srinivasan Arunachalam and Penghui Yao. Positive spectrahedra: Invariance principles and pseudorandom generators. arXiv:2101.08141, 2021.
[Bal93] Keith Ball. The reverse isoperimetric problem for Gaussian measure. Discrete Comput. Geom., 10(4):411-420, 1993.
[CDS19] Eshan Chattopadhyay, Anindya De, and Rocco A Servedio. Simple and efficient pseudorandom generators from Gaussian processes. In 34th Computational Complexity Conference (CCC 2019). Schloss Dagstuhl-Leibniz-Zentrum fuer Informatik, 2019.
[HKM13] Prahladh Harsha, Adam Klivans, and Raghu Meka. An invariance principle for polytopes. Journal of the ACM (JACM), 59(6):1-25, 2013.
[HXZ21] Yong Huang, Dongmeng Xi, and Yiming Zhao. The Minkowski problem in Gaussian probability space. Advances in Mathematics, 385:107769, 2021.
[Jam15] G. J. O. Jameson. A simple proof of Stirling's formula for the gamma function. Math. Gaz., 99(544):68-74, 2015.
[KOS08] Adam R. Klivans, Ryan O'Donnell, and Rocco A. Servedio. Learning geometric concepts via Gaussian surface area. In 49th Annual IEEE Symposium on Foundations of Computer Science, pages 541-550. IEEE, 2008.
[LM00] B. Laurent and P. Massart. Adaptive estimation of a quadratic functional by model selection. Ann. Statist., 28(5):1302-1338, 2000.
[LR10] Michel Ledoux and Brian Rider. Small Deviations for Beta Ensembles. Electronic Journal of Probability, 15:1319-1343, 2010.
[Naz03] Fedor Nazarov. On the maximal perimeter of a convex set in $\mathbb{R}^{n}$ with respect to a Gaussian measure. In Geometric aspects of functional analysis, volume 1807 of Lecture Notes in Math., pages 169-187. Springer, Berlin, 2003.
[RS15] Vijay K. Rohatgi and Ehsanes Saleh. An Introduction to probability and statistics. Wiley, 2015.
[ST17] Rocco A Servedio and Li-Yang Tan. Fooling intersections of low-weight halfspaces. In 2017 IEEE 58th Annual Symposium on Foundations of Computer Science (FOCS), pages 824-835. IEEE, 2017.
[Tao12] Terence Tao. Topics in random matrix theory, volume 132 of Graduate Studies in Mathematics. American Mathematical Society, Providence, RI, 2012.


[^0]:    *IBM T.J. Watson Research Center Srinivasan.Arunachalam@ibm.com
    ${ }^{\dagger}$ Courant Institute of Mathematical Sciences, New York University, regev@cims.nyu.edu
    ${ }^{\ddagger}$ State Key Laboratory for Novel Software Technology, Nanjing University, pyao@nju.edu.cn
    ${ }^{1} \mathrm{~A} d$-facet polytope is the special case of a spectrahedron when the matrices, $A^{(1)}, \ldots, A^{(n)}, B$ are diagonal.

